

Improving the Type System and Variable Access in the Dyninst API

Jeff Hollingsworth
John Davis



© Copyright 1999, Jeffrey K. Hollingsworth, All Rights Reserved.

Need for Type Support in Dyninst

- Access to local (stack variables)
- Complex types
 - non-integer scalars
 - structures
 - arrays
- Correctness debugging

Type Related Classes

- BPatch_type

- getName - returns the symbolic name
- getSize - returns the size of the type
- getComponents - returns the fields of struct/union
- type - returns data class (structure, union, array, ...)
- getType - return the type of the pointer, array element
- getLow, getHigh - returns bounds for arrays
- isCompatible(Bpatch_type *t2) -
test compatibility of two types

- BPatch_field

- getName - returns the field's name
- getType - returns the Bpatch_type of the field
- getOffset - returns the first byte of the field

Interaction with other Classes

- BPatch_variableExpr
 - getType - returns the type of the variable
- BPatch_image and BPatch_module
 - findType - looks up a type
- BPatch_function and BPatch_point
 - findVariable - looks up a variable in a local scope

Implementation

- Use Compiler debugger info (stab records)
 - access to user defined types
 - information about local variables
 - type information for all variables
 - line number to text segment address mapping
- Incremental parsing
 - parse stabs for a module on first use
- dyninst User can define types
 - allows the creation of new types for patched code

Stab Records May Not be Available

- Reasons for lack of Stabs
 - Programs are “stripped”
 - individual modules may not be compiled for debugging
- User type construction reduces problem
 - users can create “required” types
 - can define types for
 - global variables: often know address
 - parameters: named by position
 - define structs and array types
 - setType method of variableExpr
- Local variable access
 - not possible without stabs

Type Checking

- Ensures that snippets are type compatible
 - can disable type checking at any time
- Based on structural equivalence
 - rules:
 - scalars: same type
 - structures: each field must be compatible
 - unions: each field must be compatible
 - pointers: each points to a compatible type
 - allows more flexibility for missing types
- Error Reporting
 - snippets lack line numbers

Example of Structural Equivalence

- Patched code using a parameter struct
 - If debug info is guaranteed to be available:
 - code can access type, and refer to field
 - full type checking is possible
 - If debug info might be available:
 - can't depend on program's definition of struct
 - patch code create structure that is identical to program's version
 - permits type checking if debug info available
 - If debug info is not available
 - patch code creates structure
 - no parameter type checking possible

API Example

```
// find all variables defined in an image
BPatch_Vector<BPatch_variableExpr *> vars =
    applImage->getGlobalVariables()

for (i=0; i < vars->size(); i++) {
    BPatch_variableExpr *v = (*vars)[i];
    switch (v->getType()->type()) {
        case BPatch_scalar:
            printf("%s is a scalar of type %s\n", v->getName(),
                v->getType()->getName());
        case BPatch_structure:
            FieldVector *fields = v->getType()->getComponents();
            for (j=0; j < fields->size(); j++) {
                Bpatch_field *f = (*fields)[j];
                printf("field %s is of type %s\n", f->getName(),
                    f->getType()->getName());
            }
    }
}
```

Non-integer Scalars

- Key types

- floats - requires generating floating point expressions
- different sized integers - 16, 32, and 64 bits are needed

- Code Generation Issues

- register management
 - floats require different registers
 - 64 bit integers often need 32 bit register pairs
- expression generation
 - many instruction types needed
 - platform specific code for all supported platforms

Re-working Dyninst Code Generation

- Goals
 - support floats and ints other than 32 bit
 - enable a peephole optimizer
 - allow better register allocation
- New register abstraction
 - aware of types: int, floats, paired registers
 - allow “virtual registers” for register optimization
- Table driven instruction selection
 - eases support of multiple types
 - allows description of complex instruction
 - example: increment memory

New Dyninst Utility

- TCL-based command line tool
 - provides access to most dyninst features
 - easier to program for simple applications
 - can be used as a simple command-line debugger
 - fast conditional breakpoints
 - dynamic addition of printf's
- Command Summary
 - declare: create a new variable in the application
 - cbreak: insert conditional breakpoint
 - print: show contents of application data structures
 - at: insert a code snippet into the application
 - load, run, exit: process creation and manipulation

TCL Command Example

```
% load application
```

```
% declare int counter
```

```
% at main entry { counter = 0; }
```

```
% at importantFunc entry { counter++; }
```

```
% at main exit {
```

```
    printf("function called %d times\n", counter);
```

```
}
```

```
% run
```

Status

- Stab Parsing Working
 - currently only GNU compilers
- Array and Structure access
 - completed
- TCL Command Tool
 - mostly done - demo today
 - more features needed
- In Progress
 - local variable access
 - non-integer scalars