

# Condor on Dedicated Clusters

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The logo for Condor, featuring a large, stylized 'C' with a grey-to-black gradient and a gold outline, followed by the word 'ondor' in a gold, serif font.

# Condor on Dedicated Clusters: Overview

- Existing Implementations & Capabilities
  - NCSA Origin 2000
  - UW-Madison Compute Cluster
  - Condor Features, Tools, Tricks

# Condor on Dedicated Clusters: Overview (cont.)

## > Future Directions

- Resource Reservations
- Parallel Schedulers
- Two Scheduling Models in One Pool



# NCSA Origin 2000 and Condor

- > Massively parallel individual machines (64 to 128 CPUs each)
- > LSF scheduler used to schedule dedicated jobs on multiple CPUs
- > Condor ready to "backfill" dedicated schedule gaps with opportunistic jobs



# NCSA Origin 2000 and Condor (cont.)

- When LSF sees a scheduling gap, it can dynamically inform Condor of the number of available CPUs
  - If more nodes are free, they are advertised to a local Condor pool
  - If more nodes are needed, Condor uses job state to intelligently choose which CPUs to return to LSF



# UW-Madison's Compute Cluster

- Built from Low-Cost Commodity Parts
  - 64 Nodes
  - Each with dual-550mhz Xeon CPUs, 1GB RAM, 100Mb NIC
  - Dedicated resources
- Condor runs on all nodes unless specifically disabled for "clean-host" experiments

# UW-Madison's Compute Cluster (cont.)

- Supported by 2 quad-CPU file servers
  - 50GB storage each, 1Gb NICs to cluster
  - Dual-homed with 100Mb NIC to high-speed EMERGE research network
    - Connected to SRB, HPPS, and other data storage systems
    - Network has quality of service functionality
      - Can reserve bandwidth for transfers



# Condor & Clusters: PVM & MPI

- > PVM works well with Condor's traditional opportunistic model
  - Can dynamically adjust to # of available nodes
- > MPI doesn't work as well
  - Requires fixed number of nodes
  - More suited to dedicated resources





# Condor & Clusters: Features & Tools

## > Parallel checkpointing

- Problem:

- Traditional Condor pools often can't assume each machine has adequate disk/network to checkpoint jobs
- Central checkpoint server is a *potential* bottleneck



# Parallel Checkpointing (cont.)

- Observation:
  - Clusters often have adequate resources
    - Great network between nodes
    - Lots of local disk
- Solution:
  - Checkpoint to multiple servers
  - Servers run on cluster nodes themselves
  - Scales very well: You can have a checkpoint server on every node!



# Future Directions

- Reservations
- Parallel Scheduling
- Dual Scheduling Across 1 Pool
  - Dedicated scheduling
  - Opportunistic scheduling



# Future Directions: Reservations

- Two kinds of reservations: interactive users & future jobs
- Use by resource owner not an issue
- Guaranteed reservations needed
- Need two more entities in Condor:
  - Reservation manager
  - Reservation enforcer



# Reservations (cont.)

- Reservation Manager:
  - Reservations are represented as a ClassAd -- can use all the existing technology:
    - Persistent storage
    - Network communication layer
    - Visualization



# Reservations (cont.)

- Reservation Manager (cont.):
  - Support requests by jobs in the system
  - Supports a GUI for interactive users
- Reservation Enforcer:
  - Uses ClassAd matchmaking technology
  - Vacate nodes in advance to avoid flooding the network
  - Current plan: use the Eventd

# Future Directions: Parallel Scheduling

- Co-scheduling of multiple hosts
  - MPI Job ClassAd might require  $N$  nodes
    - You want all or nothing, or you can have deadlock
  - Other jobs might require co-scheduling:
    - A multi-threaded application might want to claim multiple CPUs on a single SMP machine
  - Requires "gang-matching"

# Future Directions: Dual Scheduling Across One Pool

- Condor's hierarchical and parallel scheduling architecture will enable dedicated and opportunistic schedulers to coexist and efficiently share dedicated and non-dedicated resources alike
- Resources will be able to migrate between scheduling systems





# Dual Scheduling Across One Pool (cont.)

- During schedule gaps, dedicated compute machines will become available to the opportunistic scheduler
- However, dedicated machines will always "prefer" the dedicated scheduler and will return as soon as they are needed

Questions  
and  
Thank You!