

# Using Automated Performance Modeling to Find Scalability Bugs in Complex Codes



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# **Analytical performance modeling**



### Disadvantages

- Time consuming
- Danger of overlooking unscalable code





Generate an empirical model for each part of the program automatically

- Run a manageable number of small-scale performance experiments
- Launch our tool
- Compare extrapolated performance to expectations

#### Key ideas

- Exploit that space of function classes underlying such model is small enough to be searched by a computer program
- Abandon model accuracy as the primary success metric and rather focus on the binary notion of scalability bugs
- Create requirements models alongside execution-time models



**Scalability bug detector** 

### Input

- Set of performance measurements (profiles) on different processor counts
  - $\{p_1, ..., p_{max}\}$  w/ weak scaling
- Individual measurement broken down by program region (call path)

### Output

- List of program regions (kernels) ranked by their predicted execution time at target scale p<sub>t</sub> > p<sub>max</sub>
- Or ranked by growth function  $(p_t \rightarrow \infty)$

- Not 100% accurate but good enough to draw attention to right kernels
- False negatives when phenomenon at scale is not captured in data
- False positives possible but unlikely
- Can also model parameters other than p



### Workflow





# **Model generation**

### **Performance Model Normal Form (PMNF)**

$$f(p) = \sum_{k=1}^{n} c_k \cdot p^{i_k} \cdot \log_2^{j_k}(p)$$

- Not exhaustive but works in most practical scenarios
- An assignment of *n*,  $i_k$  and  $j_k$  is called model hypothesis
- $i_k$  and  $j_k$  are chosen from sets  $I, J \subset \mathbf{Q}$
- *n*, |*I*|, |*J*| don't have to be arbitrarily large to achieve good fit

Instead of deriving model through reasoning, make reasonable choices for n, I, J and try all assignment options one by one

Select winner through cross-validation



# **Model refinement**



- Start with coarse approximation
- Refine to the point of statistical shrinkage
- $\rightarrow$  Protection against over-fitting



## **Requirements modeling**



Disagreement may be indicative of wait states







We demonstrate that our tool

- identifies a scalability issue in a code that is known to have one
- does not identify a scalability issue in a code that is known to have none
- identifies two scalability issues in a code that was thought to have only one

Test platform: IBM Blue Gene/Q Juqueen in Jülich

$$I = \{\frac{0}{2}, \frac{1}{2}, \frac{2}{2}, \frac{3}{2}, \frac{4}{2}, \frac{5}{2}, \frac{6}{2}\}$$
$$J = \{0, 1, 2\}$$
$$n = 5$$



## Sweep3D

Solves neutron transport problem

- 3D domain mapped onto 2D process grid
- Parallelism achieved through pipelined wave-front process



$$t^{comm} = [2(p_x + p_y - 2) + 4(n_{sweep} - 1)] \cdot t_{msg}$$
$$t^{comm} = c \cdot \sqrt{p}$$







Kernel	Runtime[%] p <sub>t</sub> =262k	Increase <u>t(p=262k)</u> t(p=64)	Model [s] t = f(p) p <sub>i</sub> ≤ 8k		Predictive error [%] p <sub>t</sub> =262k	
sweep → MPI_Recv	65.4	16.5	4.0√p			5.1
sweep	20.9	0.2	582.2	#bytes = co #msg = co	onst. nst.	0.01
global_int_sum → MPI_Allreduce	12.9	18.7	1.1√p + 0.03√p∙log	(p)		13.6
sweep $\rightarrow$ MPI_Send	0.4	0.2	11.5 + 0.1	p∙log(p)		15.4
source	0.3	0.04	6.7 + 9.1•1	0⁻⁵log(p)		0.01



Sweep3D (3)







MILC/su3\_rmd – code from MILC suite of QCD codes with performance model manually created by Hoefler et al.

 Time per process should remain constant except for a rather small logarithmic term caused by global convergence checks

Kernel	Model [s] t=f(p) p <sub>i</sub> ≤ 16k	Predictive Error [%] p <sub>t</sub> =64k
compute_gen_staple_field	0.02	0.4
g_vecdoublesum $\rightarrow$ MPI_Allreduce	6.3 •10 <sup>-6</sup> •log <sup>2</sup> p	0.01
mult_adj_su3_fieldlink_lathwec	0.004	0.04





Core of the Community Atmospheric Model (CAM)

 Spectral element dynamical core on a cubed sphere grid



Kernel	Model [s] t = f(p) p <sub>i</sub> ≤15k	Predictive error [%] p <sub>t</sub> = 130k
Box_rearrange->MPI_Reduce	$0.03 + 2.5 \cdot 10^{-6} p^{3/2} + 1.2 \cdot 10^{-12} p^3$	57.0
Vlaplace_sphere_vk	49.5	99.3
Compute_and_apply_rhs	48.7	1.7





Core of the Community Atmospheric Model (CAM)

 Spectral element dynamical core on a cubed sphere grid



Kernel	Model [s] t = f(p) p <sub>i</sub> ≤43k	Predictive error [%] p <sub>t</sub> = 130k
Box_rearrange->MPI_Reduce	3.6•10 <sup>-6</sup> p <sup>3/2</sup> + 7.2•10 <sup>-13</sup> p <sup>3</sup>	30.3
Vlaplace_sphere_vk	24.4 + 2.3•10 <sup>-7</sup> p <sup>2</sup>	4.3
Compute_and_apply_rhs	49.1	0.8





#### Two issues

Number of iterations inside a subroutine grows with p<sup>2</sup>

- Ceiling for up to and including 15k
- Developers were aware of this issue and had developed work-around

Growth of time spent in reduce function grows with p<sup>3</sup>

- Previously unknown
- Function invoked during initialization to funnel data to dedicated I/O processes
- Execution time at 183k ~ 2h, predictive error ~40%

ANR DFG

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Automated performance modeling is feasible

Generated models accurate enough to identify scalability bugs or show their absence with high probability

Advantages of mass production also performance models

- Approximate models are acceptable as long as the effort to create them is low and they do not mislead the user
- Code coverage is as important as model accuracy

### Future work

- Study influence of further hardware parameters
- More efficient traversal of search space (allows more model parameters)
- Integration into Scalasca



## Acknowledgement





