



German Research School  
for Simulation Sciences

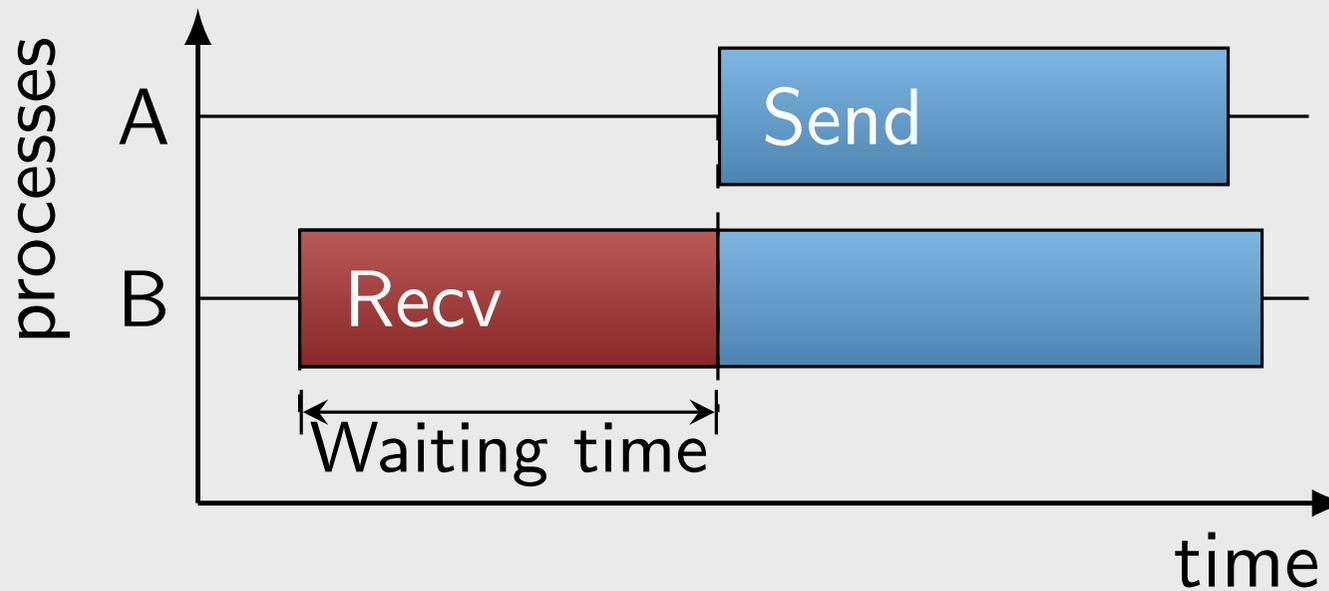
## Catching Idlers with Ease: A Lightweight Wait-State Profiler for MPI Programs



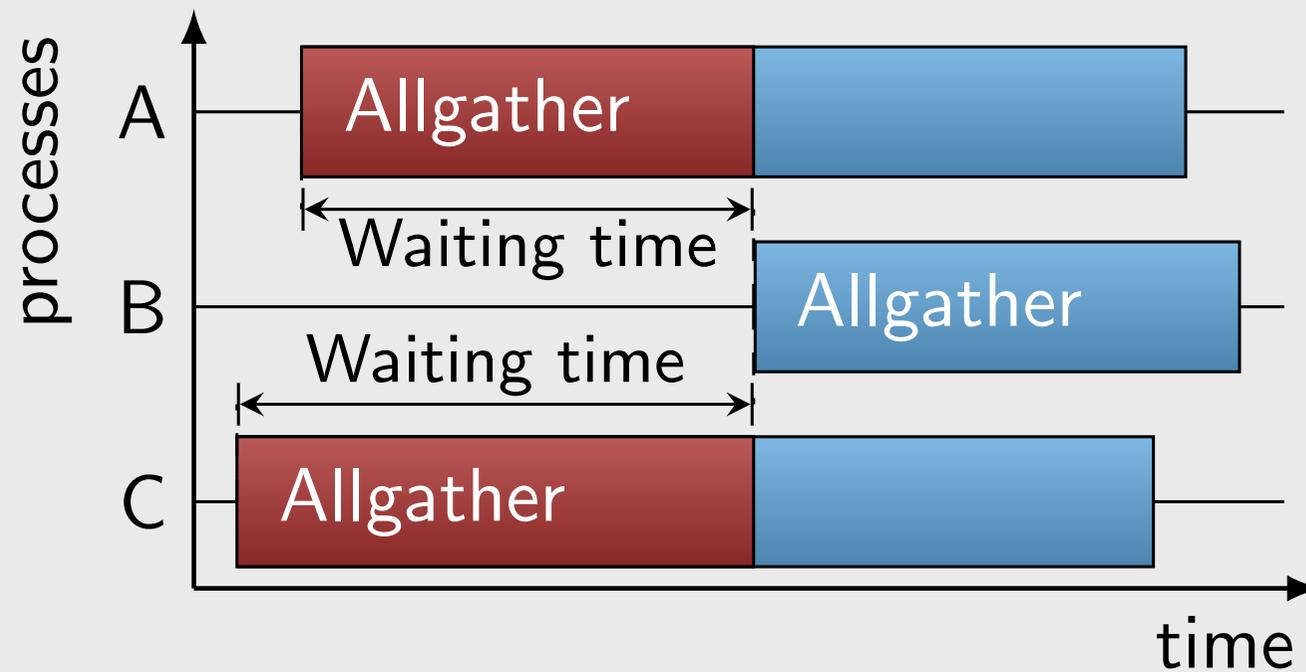
Guoyong Mao, David Böhme, Markus Geimer, Marc-André Hermanns,  
Daniel Lorenz and Felix Wolf

Petascale Tools Workshop, Madison, WI, USA, August 4, 2014

## Late sender



## Wait an NxN



## What we want to know



## What we measure



Execution time



Execution time



Execution time

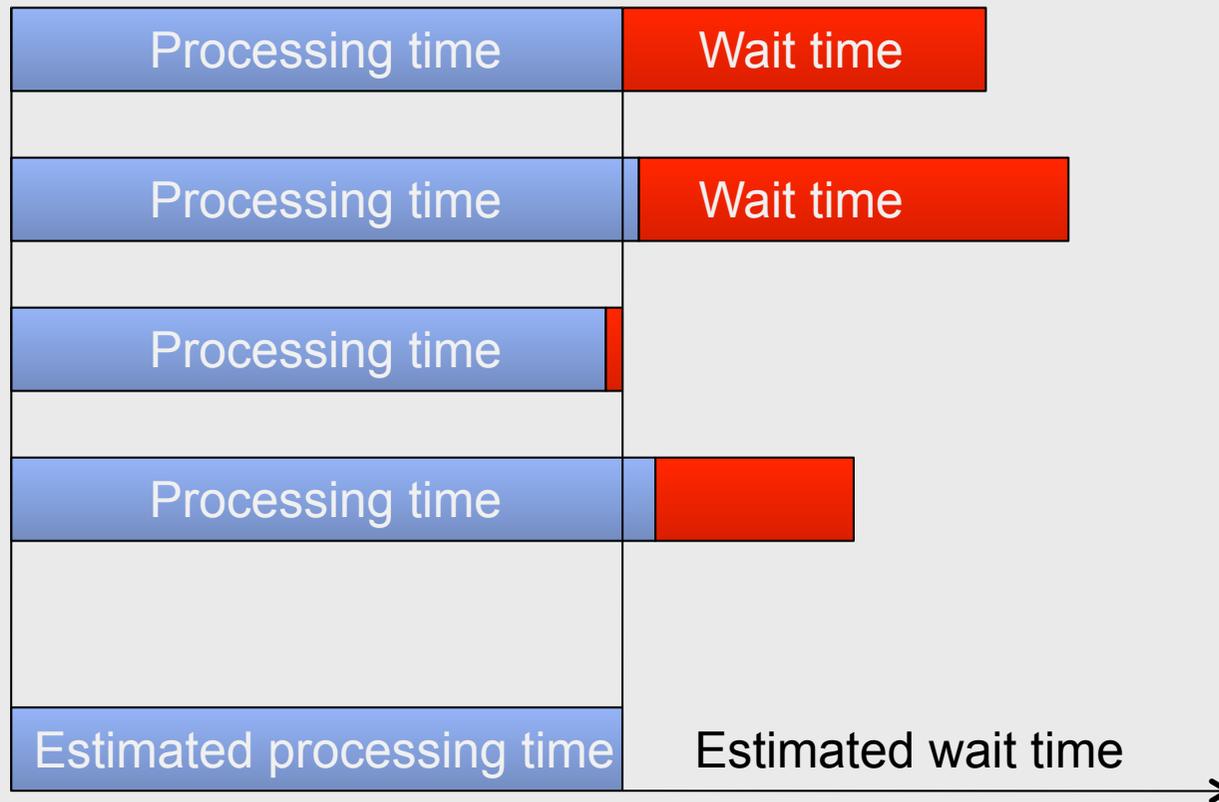


Execution time

## The minimum idea



## The minimum idea





## Considered parameters

- We consider
  - MPI function
  - Message size
  - Receiver rank
- Other possible parameters
  - Sender rank
  - Data type
- Tradeoff between
  - Number of samples for a meaningful minimum and amount data
  - Parameters considered
- Need to find the relevant parameters.

## Algorithm

For every combination of

- MPI function
- Message size class
- Process

record the

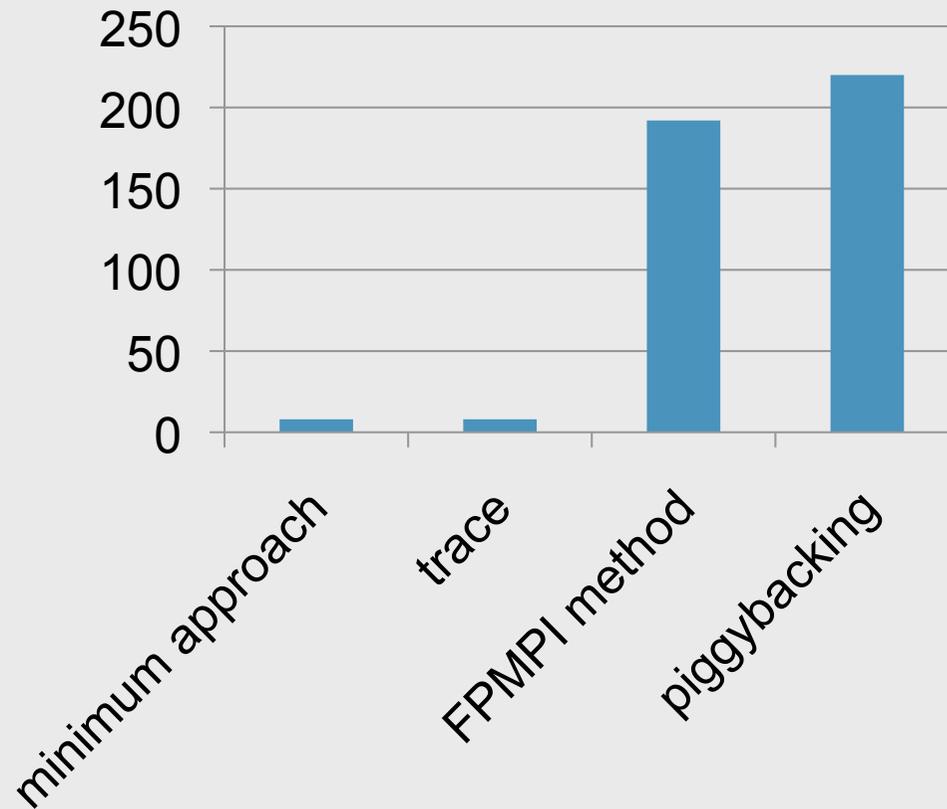
- Minimum execution time

For every combination of MPI call path and message size class record the

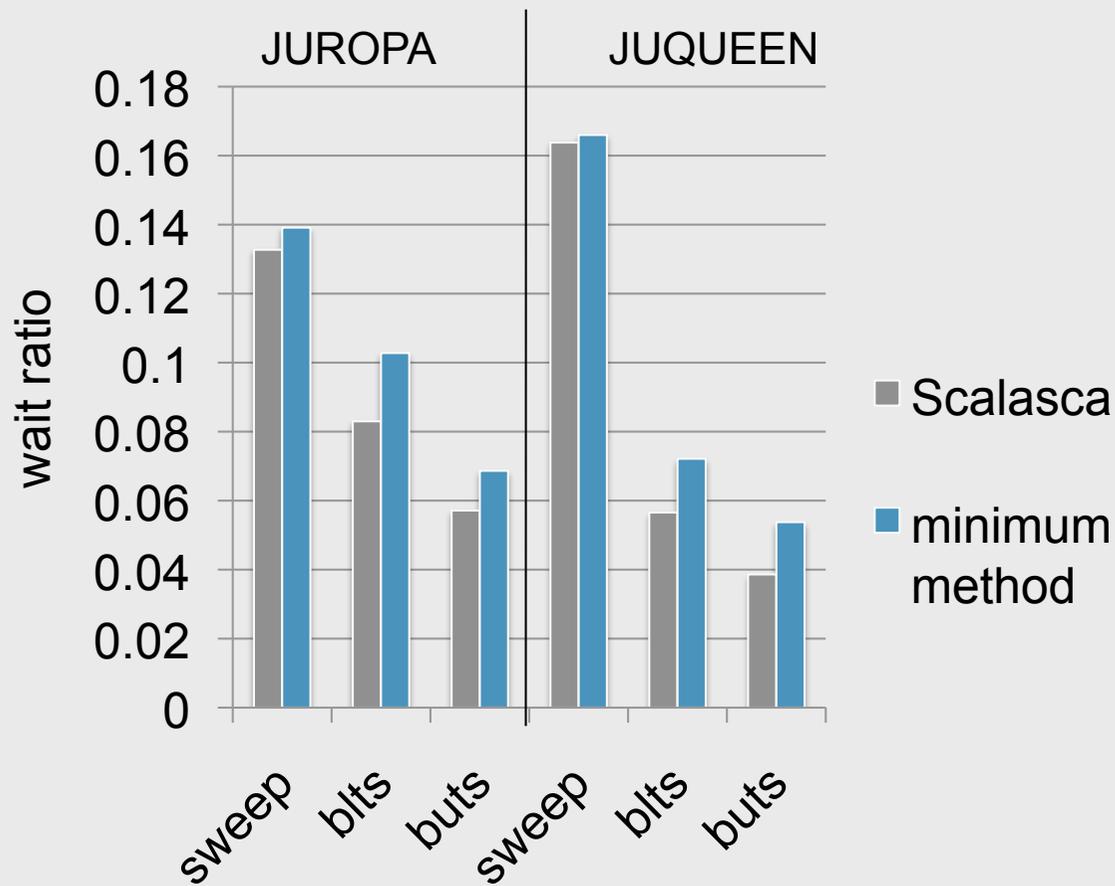
- Number of visits
- Total execution time

At the end of the profiling run, subtract the minimum from the execution time for every visit to calculate the wait time.

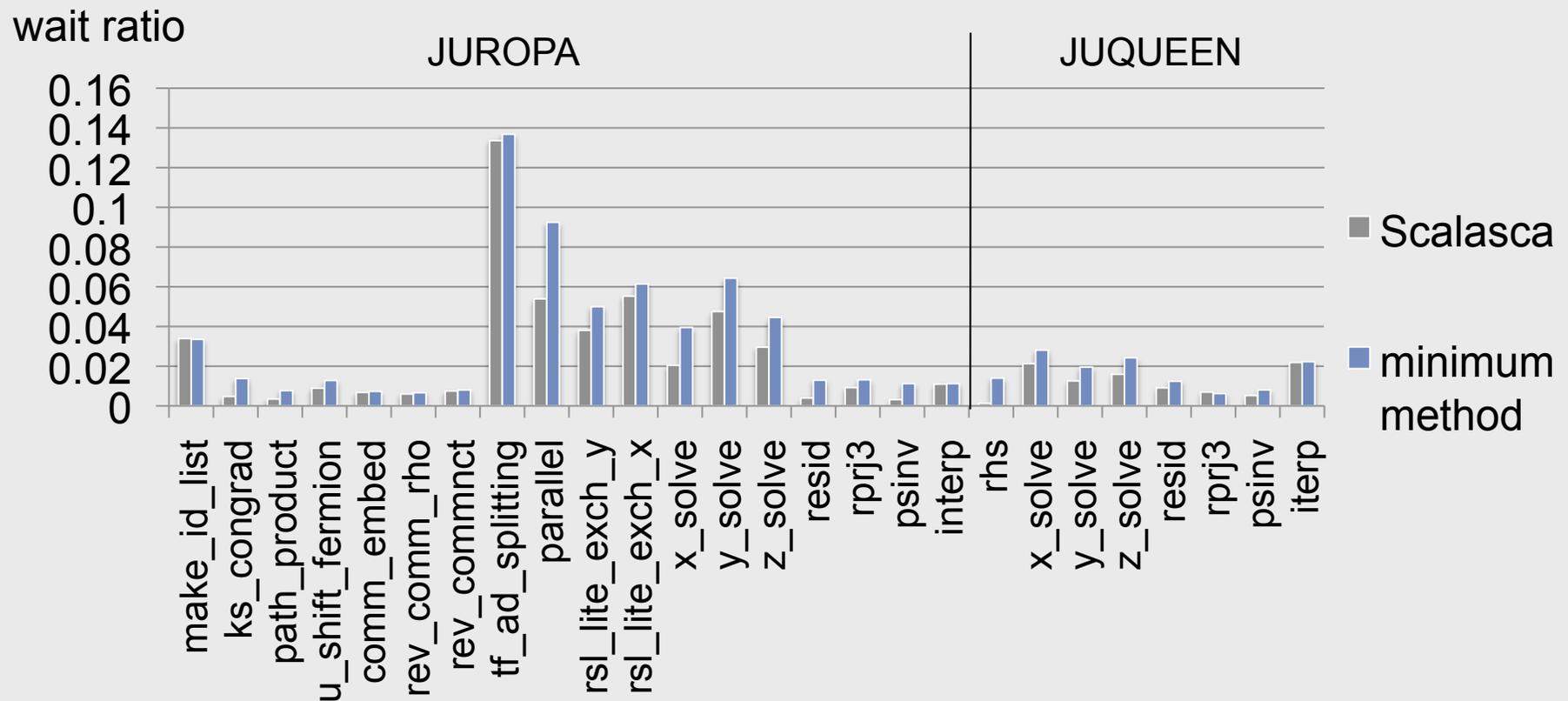
## Per-call overhead increase compared to profiling overhead w/o wait state analysis (%)



## Accuracy MPI\_Recv



## Accuracy MPI\_Wait



## Accuracy MPI\_Wait

wait ratio

0.16  
0.14  
0.12  
0.1  
0.08  
0.06  
0.04  
0.02  
0

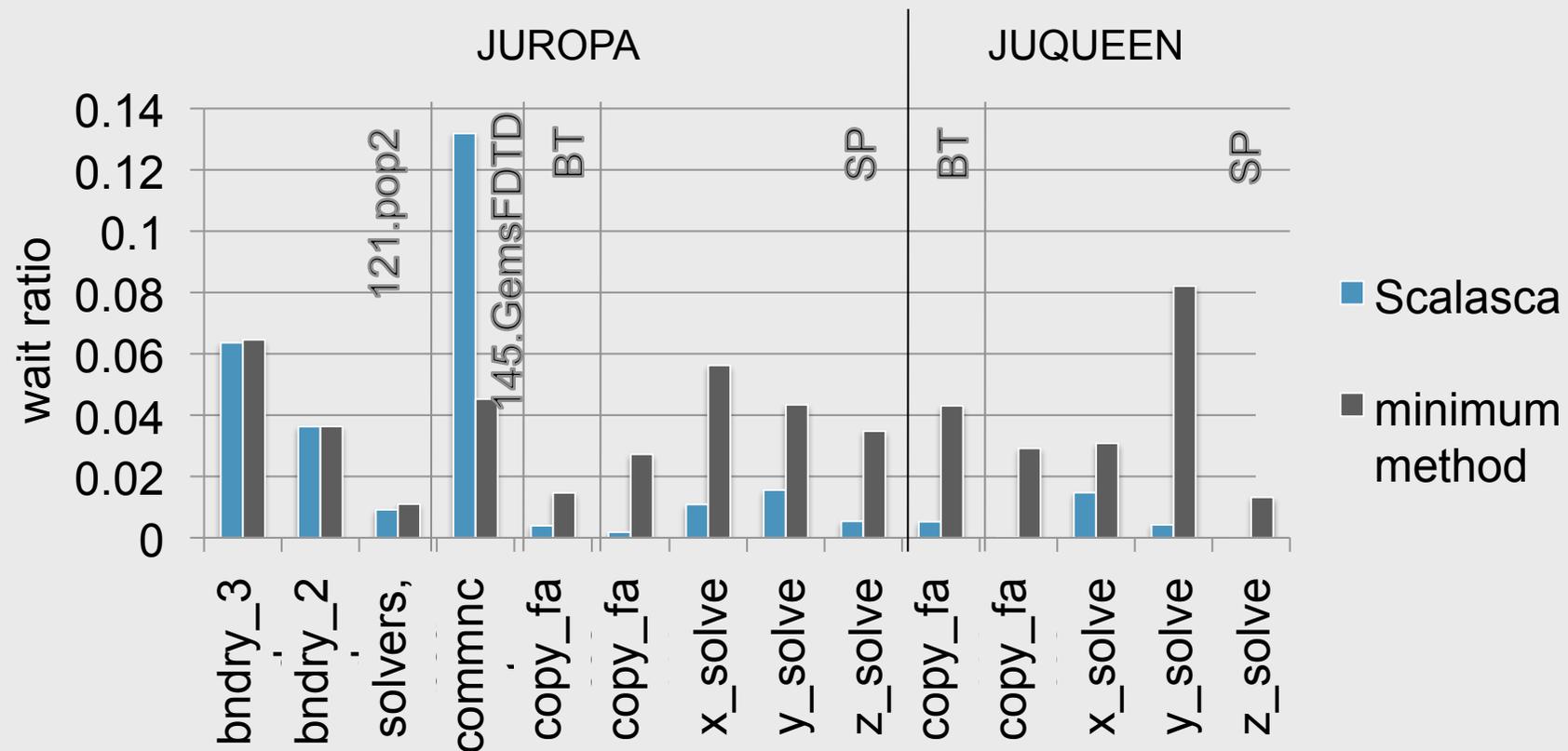
JUROPA

JUQUEEN

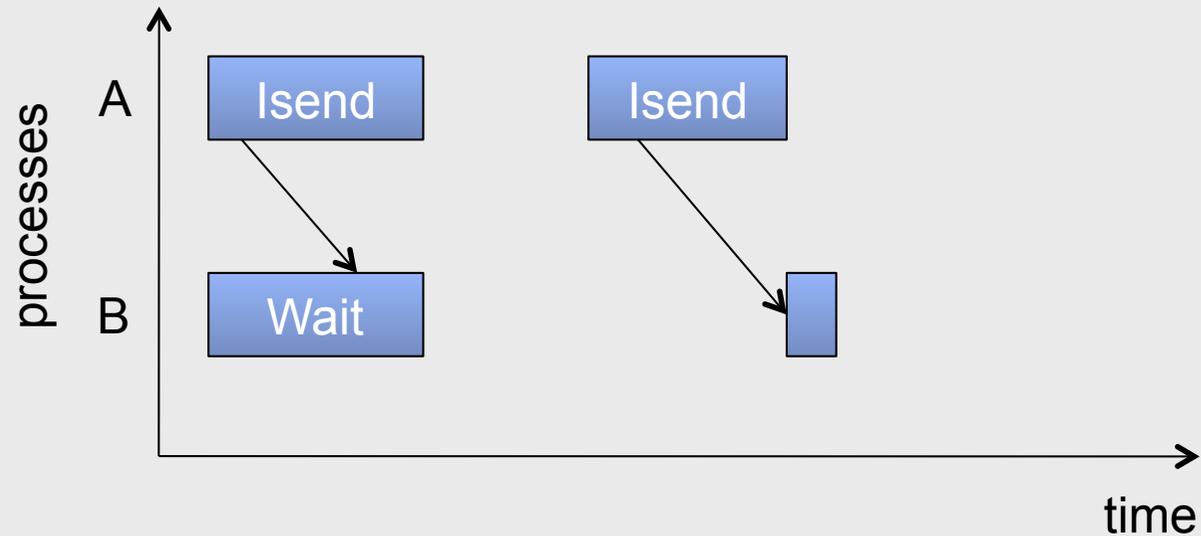
make\_id\_list  
ks\_congrad  
path\_product  
u\_shift\_fermion  
comm\_embed  
rev\_comm\_rho  
rev\_commct  
tf\_ad\_splitting  
parallel  
rsl\_lite\_exch\_y  
rsl\_lite\_exch\_x  
x\_solve  
y\_solve  
z\_solve  
resid  
rprj3  
psinv  
interp  
rhs  
x\_solve  
y\_solve  
z\_solve  
resid  
rprj3  
psinv  
interp

■ Scalasca  
■ minimum method

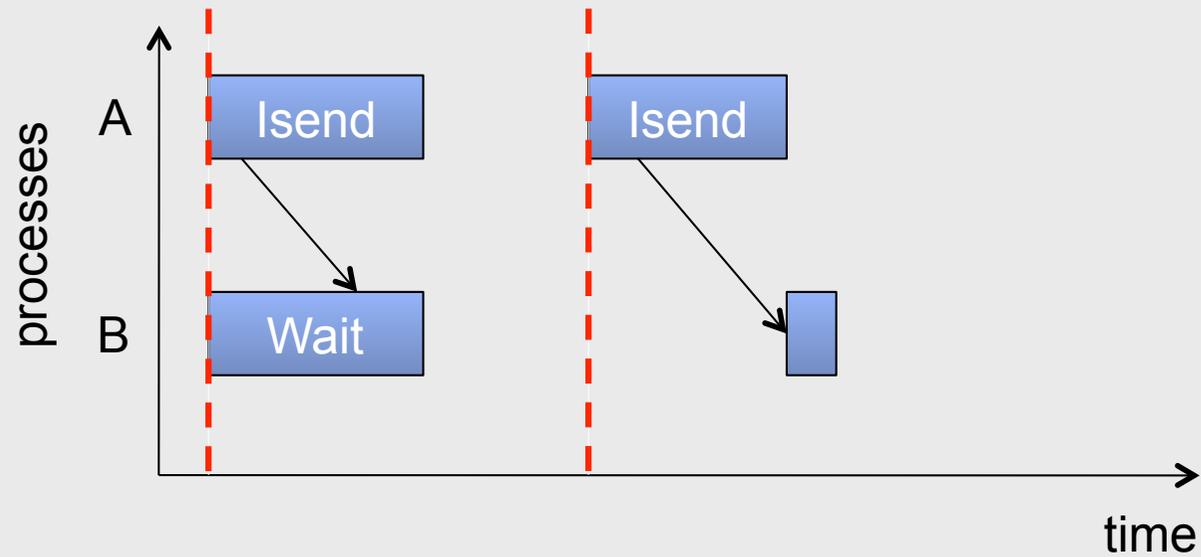
## Accuracy MPI\_Waitall



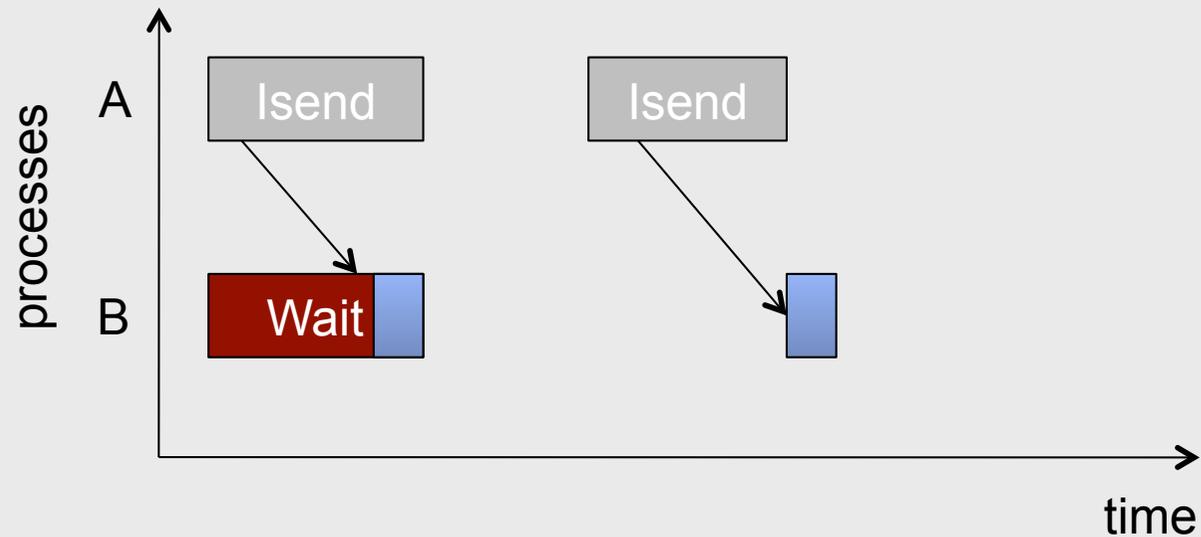
## Non-blocking communication



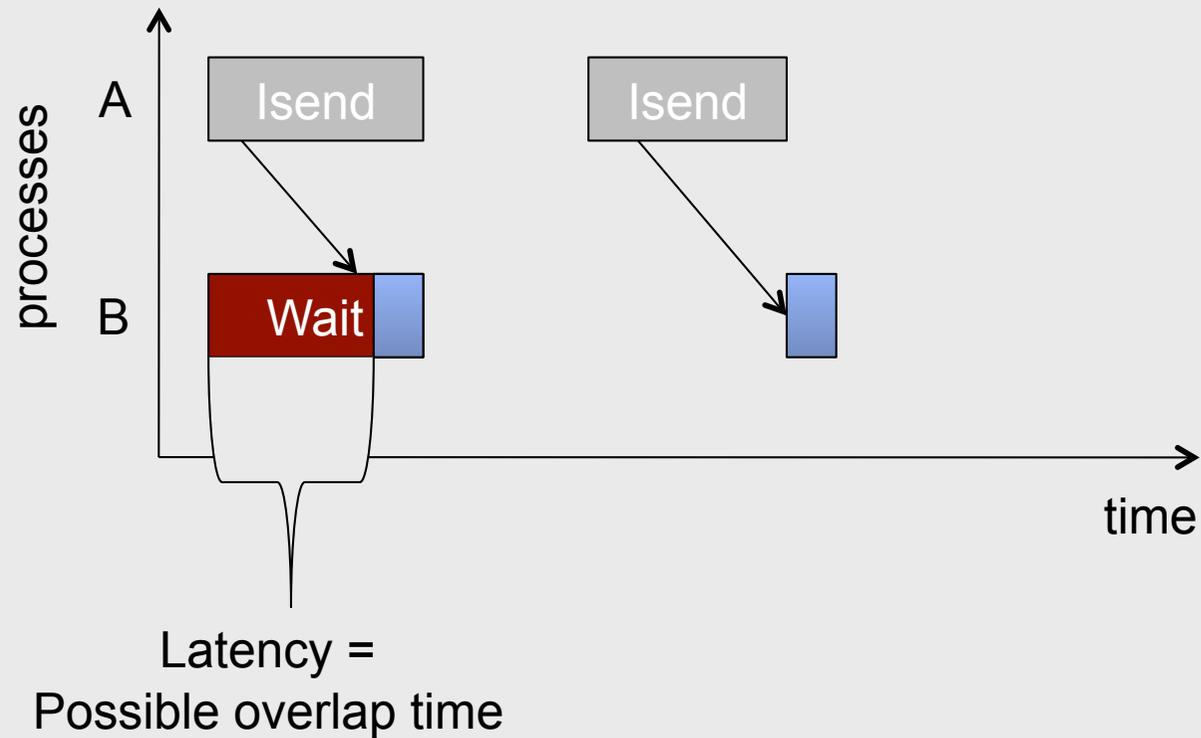
## Scalasca detects no wait state



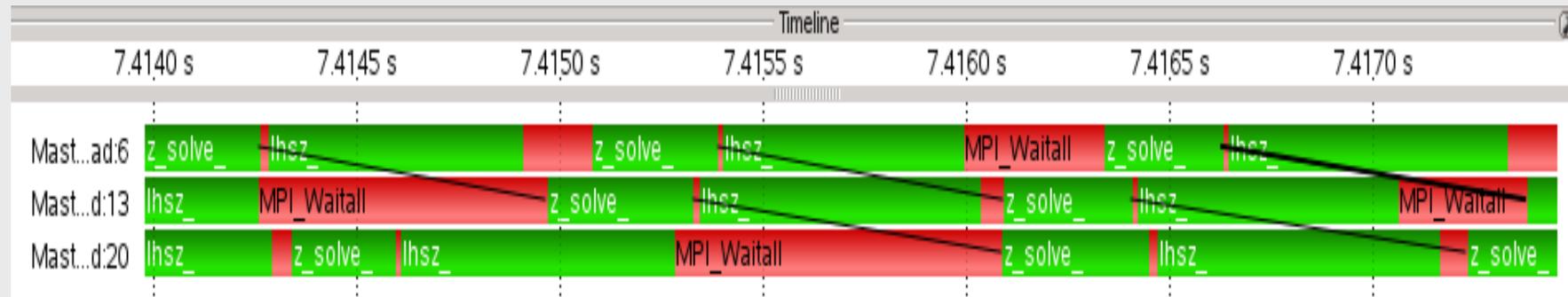
## Minimum approach does calculate wait states



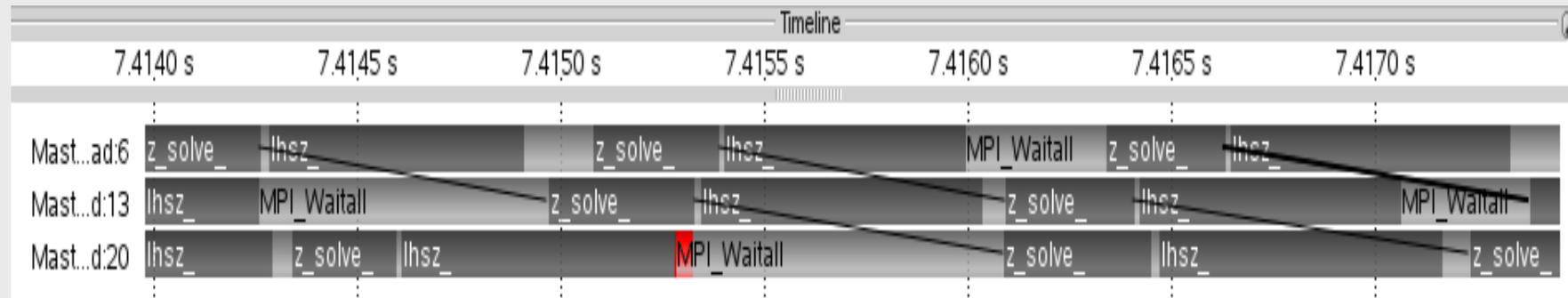
## But is this wrong for performance analysis?



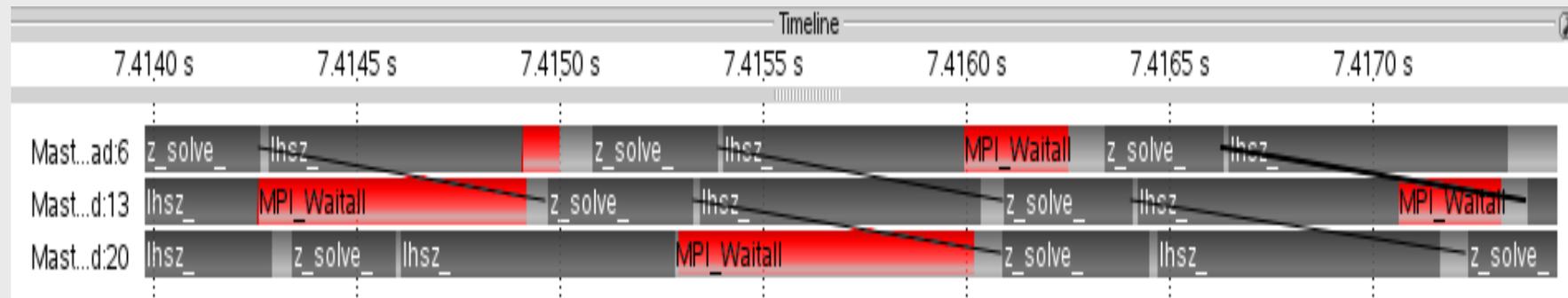
## Detailed example from SP



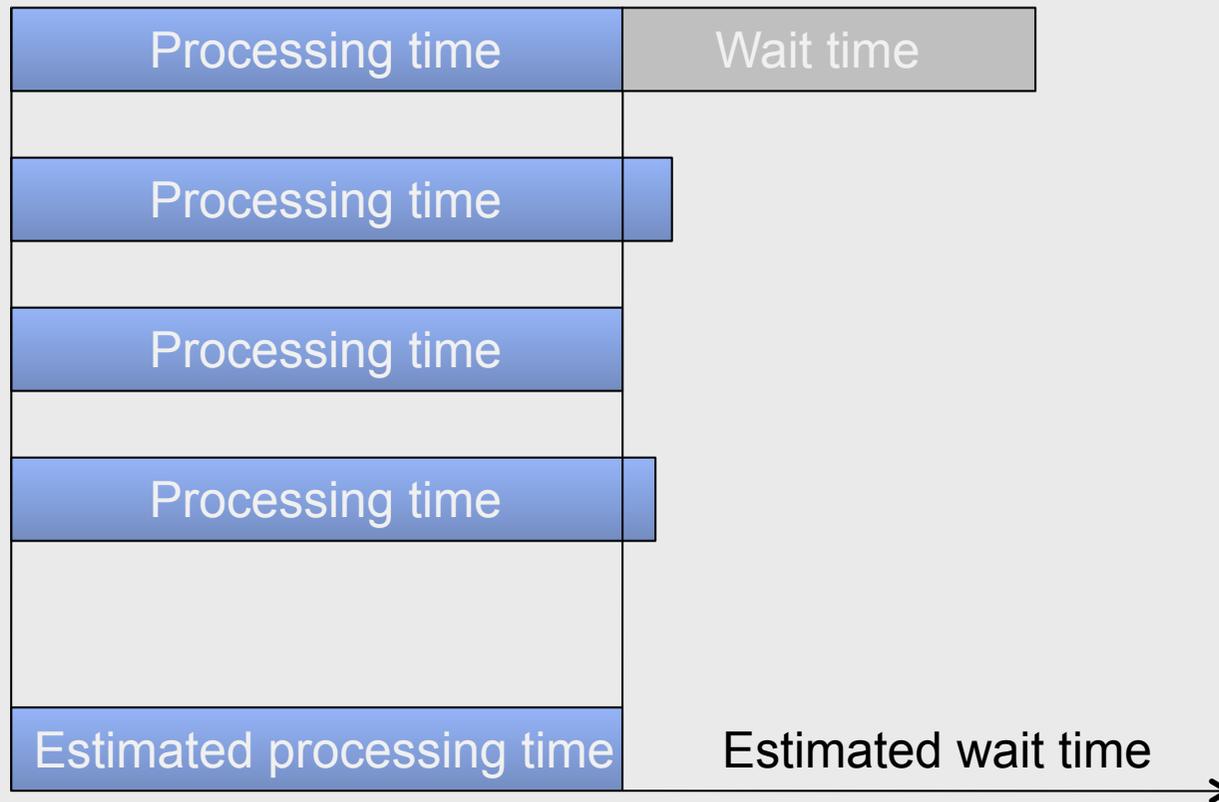
## Wait time according to Scalasca



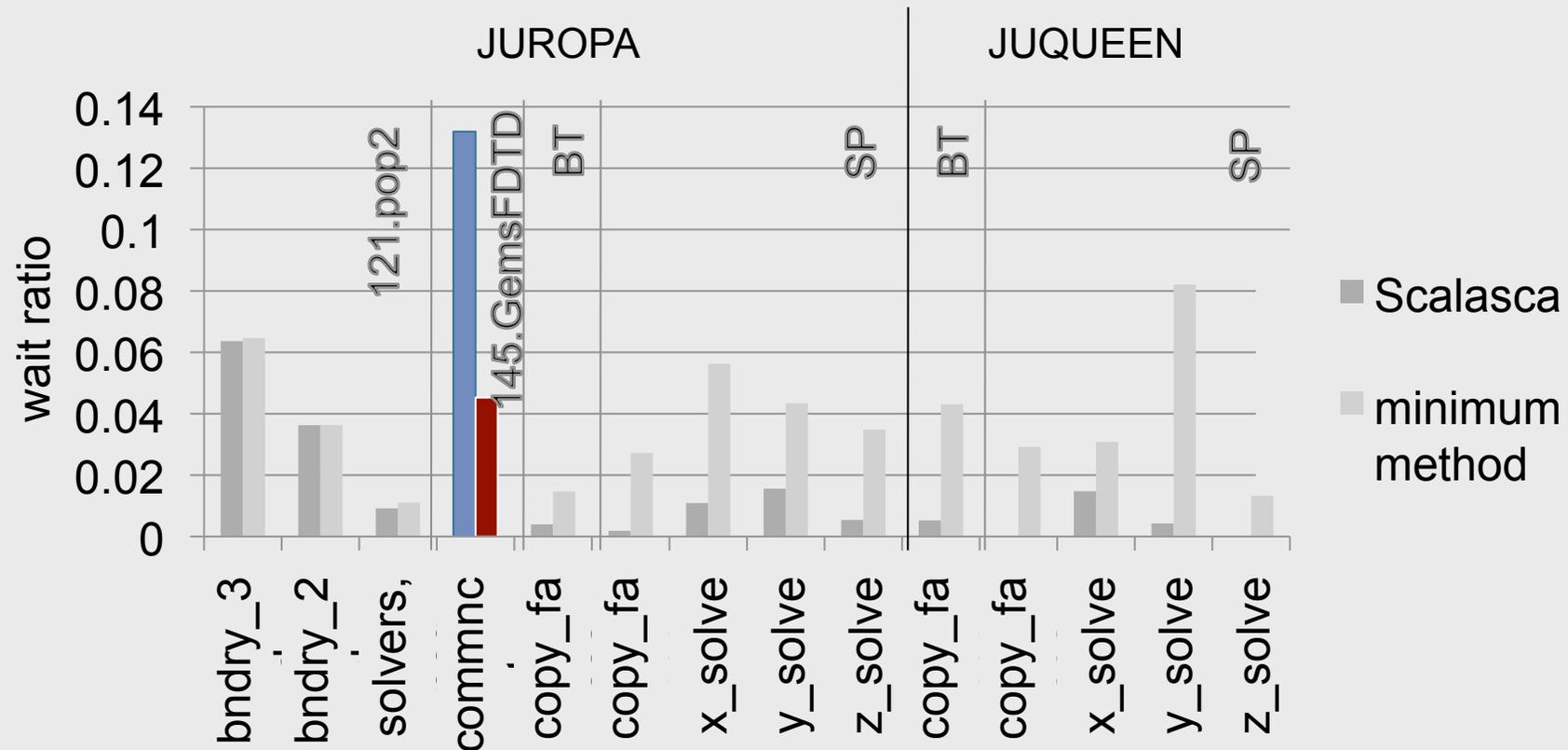
## Wait time according to minimum method



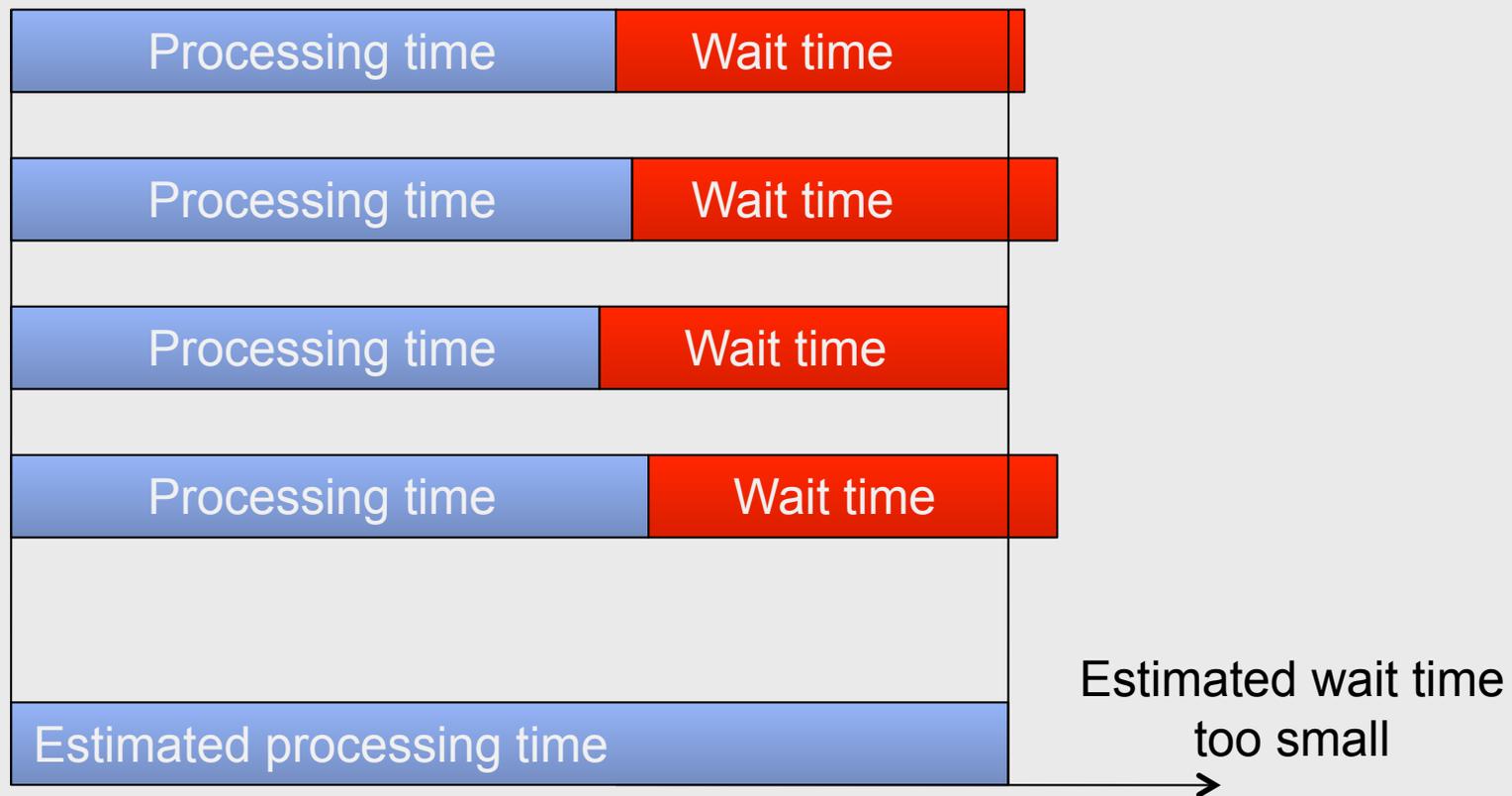
## Jitter may cause a little higher wait time



## Accuracy MPI\_Waitall



## Static imbalance

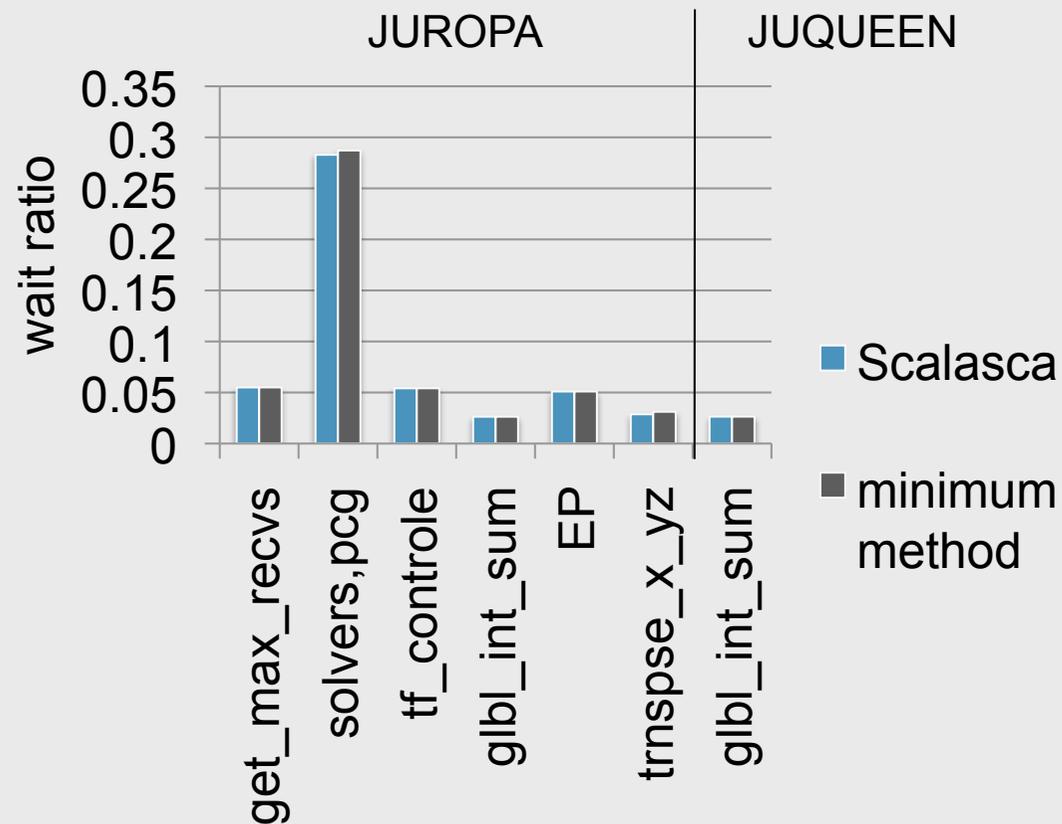




## Static imbalance

- Calculating global minima could resolve process local static imbalances
  - Reduction operation after measurement
  - No dilation at measurement time
- Loose sender/receiver parameterization of minima
- For collective operations, global minima were better

## Accuracy for Wait at NxN



## Conclusion (1)

- Minimum method works for the estimation of blocking and non-blocking communication
  - For blocking communication results similar to Scalasca
  - For non-blocking communication, in Waitall wait time do not match the Scalasca analysis.
- Low runtime overhead
- No trace recording or piggybacking
- May not produce 100% accurate numbers, but
  - Sufficient accuracy to locate performance problems
  - Point to places where we might want to investigate further with trace analysis

## Conclusion (2)

- Detection of good minimum crucial
  - Static imbalance
  - Tradeoff between number of parameters and number of samples
- Jitter may lead to minor increase of measured wait state
- For non-blocking communication
  - Count possible overlap time
  - Might be larger than pure Late Sender time
  - Isn't this even more accurate to estimate the optimization potential?



## Reference

Guoyong Mao, David Böhme, Marc-André Hermanns, Markus Geimer, Daniel Lorenz, Felix Wolf: *Catching Idlers With Ease: A Lightweight Wait-State Profiler for MPI Programs*. In: EuroMPI '14: Proc. Of the 21<sup>st</sup> European MPI Users' Group Meeting, Tokyo, Japan, Sep. 9-12, 2014